

# Template Attacks on ECDSA on a 32-bit ARM

## Theory

### ECDSA Signature Generation

1.  $e = \text{HASH}(m)$
2. random integer  $k$  from  $[1, n - 1]$
3.  $r = x(\text{mod } n)$ , where  $(x, y) = kG$ .
4. Calculate  $s = k^{-1}(e + rd_A)(\text{mod } n)$ .
5. The signature is the pair  $(r, s)$ .

### Template Based SPA Attack

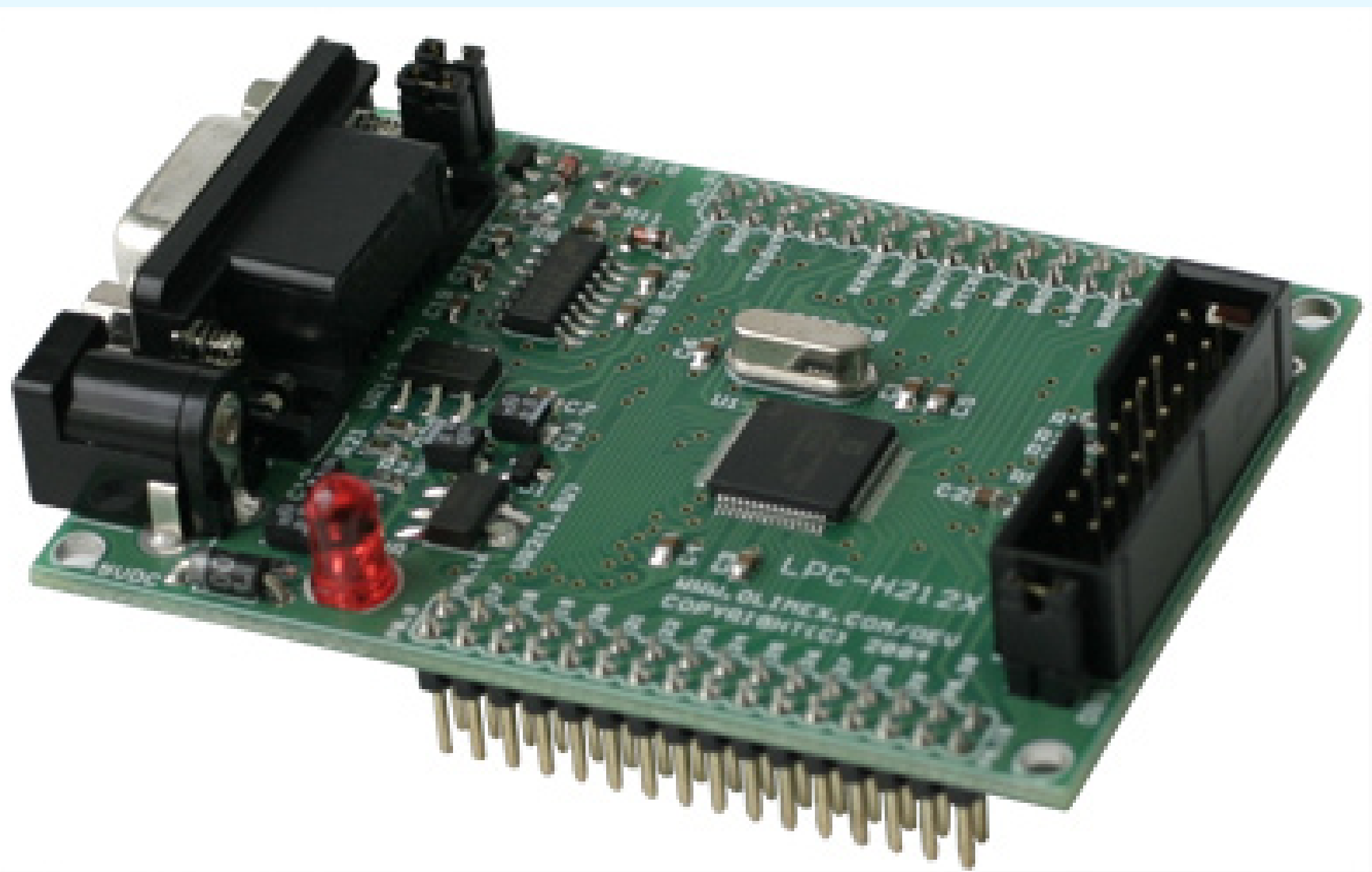
- Superior over standard SPA, succeeds even for SPA resistant algorithms
- Describes power consumption with a probability distribution
- Allows classification of all kinds of instructions, operations or features
- Single shot attack

### Attack on ECDSA

- ECDSA is well suited for template attacks
- Basepoint  $G$  is known and intermediate curve points can be precomputed and classified for a guessed key  $k$
- Successful attacks on the ephemeral key  $k$  reveal the secret key  $d_A$

## ECC Implementation

- Optimized for 32-bit ARM7 platform, which is used in many modern embedded systems
- NIST curve P192
- Binary and windowed double and always add method (both attacked successfully)
- Standard and constant runtime GF(p) arithmetic (both attacked successfully)

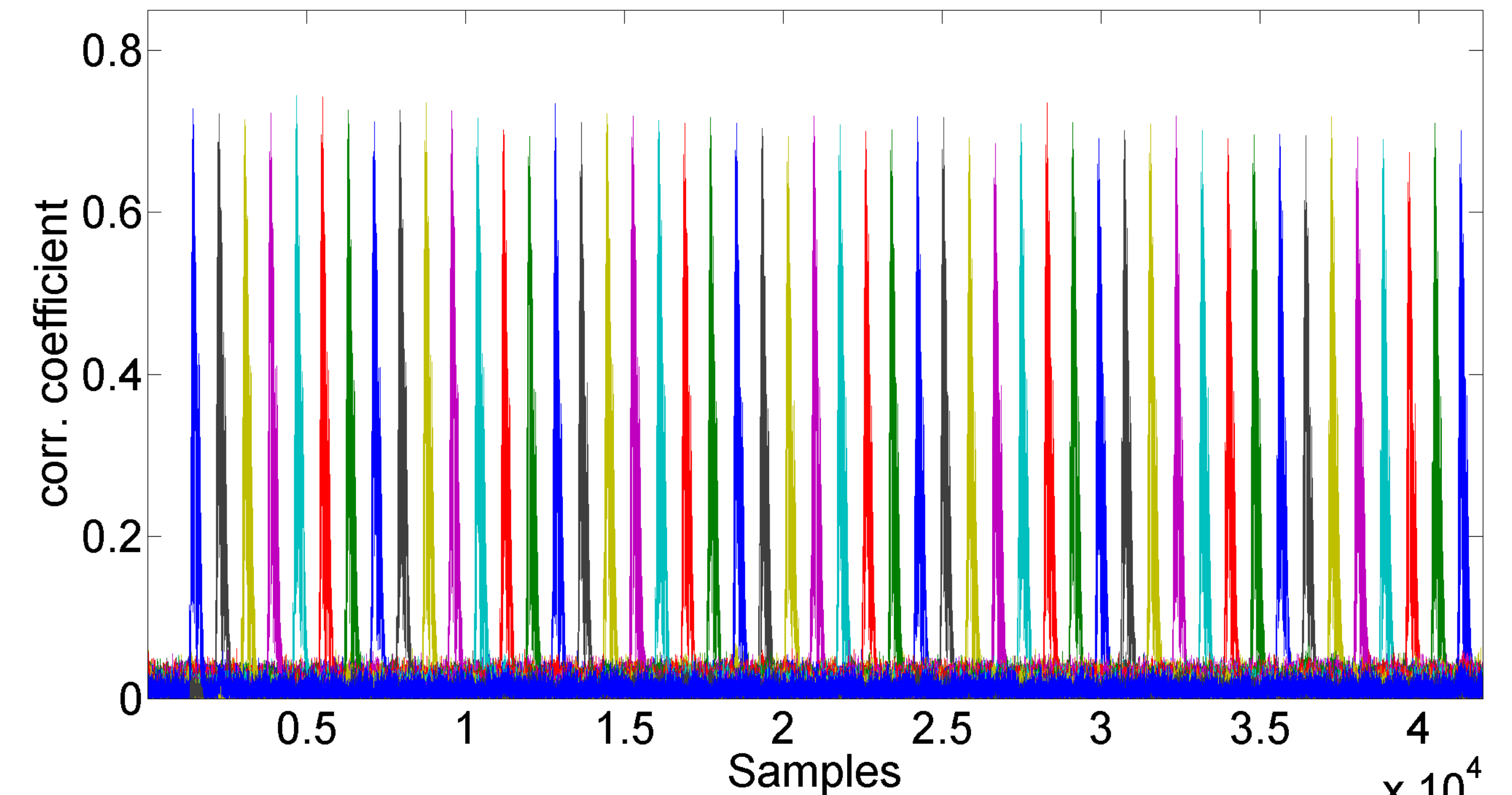


All implementations have been attacked on this ARM7 based LPC2124 32-bit microcontroller

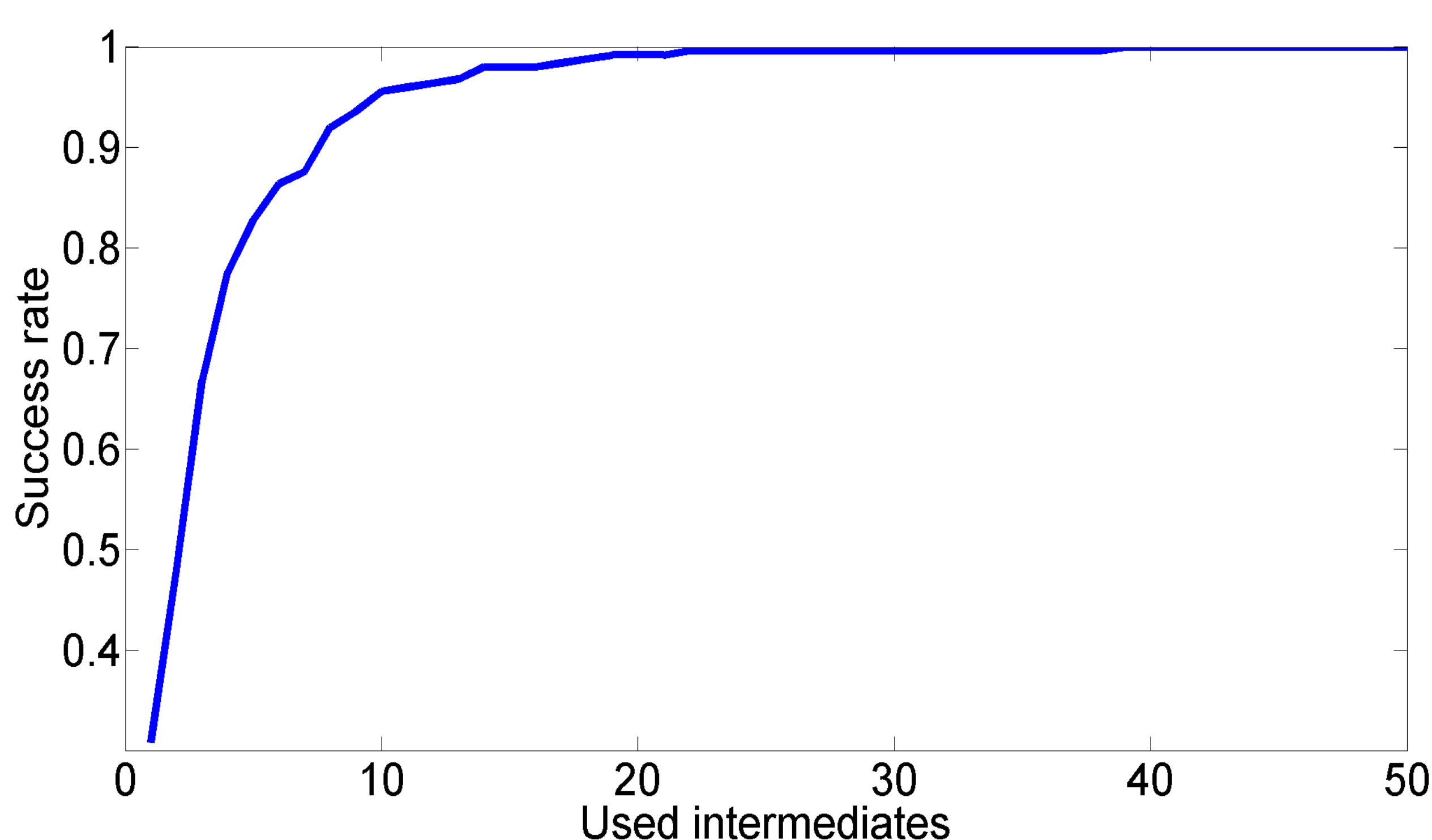
## Practical Attack

### Used Leakage

- Standard GF(p) implementations leak dramatically due to key-dependent reductions
- Constant runtime GF(p) arithmetic has DPA leakage



For the P192 curve a single group operation provides over 100 intermediate values, larger curves even more



If the template performance for one intermediate is larger than 50% it is possible to constructively combine them. Considering about 40 intermediate values is enough to achieve a success rate beyond 99.99%.

### Templates for Several Curve Points

- Allows optimal template point selection
- Best DPA point per intermediate
- Use only intermediates with large Hamming distance
- Leads to a success rate beyond 99.99%